Discrete Mathematics - Problem Session -2 - All about FOL - Solutions.

- 1. We cannot help everyone, but everyone can help some one. $\forall x \exists y(\neg help(x, y)) \land \forall x \exists y(help(x, y))$ $\forall x(\neg \forall y(help(x, y))) \land \forall x \exists y(help(x, y))$
- 2. No one who has no complete knowledge of himself will ever have a true understanding of another. This statement means, there exists one who has no knowledge of himself will ever have a true understanding of another is false. $\neg[\exists x(\neg know(x, x) \land \exists y(understand(x, y)))]$

 $\forall x (\neg know(x, x) \rightarrow \forall y (\neg understand(x, y)))$

- 3. Negate the following: $\forall x \exists \epsilon ((x > 0 \land \epsilon > 0) \land \forall y(y > 0 \rightarrow x y \ge \epsilon))$ solution: $\exists x \forall \epsilon ((x > 0 \land \epsilon > 0) \rightarrow \exists y(y > 0 \land x - y < \epsilon))$ NOTE: while applying negation, \forall flips to \exists and vice versa. Similarly, \rightarrow flips to \land and vice versa.
- 4. Some Republicans like all Democrats. No Republican likes any Socialist. Therefore, no Democrat is a Socialist. $\exists x(R(x) \land \forall y(D(y) \rightarrow like(x, y))) (1)$ $\neg(\exists x(R(x) \land \exists y(S(y) \land like(x, y))) \equiv \forall x(R(x) \rightarrow \forall y(S(y) \rightarrow \neg like(x, y))) - (2)$ $\forall x(D(x) \rightarrow \neg S(x)). - (3)$

The above claim is true, we shall present a direct proof. (3) EI,UI of (1) $R(a) \land (D(b) \rightarrow like(a, b))$, for some a and any b(4) UI of (2) $R(a) \rightarrow (S(b) \rightarrow \neg like(a, b))$, for any a and bNote: $P \land Q \rightarrow P$ is a tautology. (5) From (3), R(a)Note: $P \land (P \rightarrow Q)$ is a tautology. (6) From (5) and (4), $(S(b) \rightarrow \neg like(a, b))$ (7) Contrapositive of (6); $like(a, b) \rightarrow \neg S(b)$ Note: $((P \rightarrow Q) \land (Q \rightarrow R)) \rightarrow (P \rightarrow R)$ is a tautology. (8) From (3), we get $(D(b) \rightarrow like(a, b))$ (9) From (8) and (7), we get $D(b) \rightarrow \neg S(b)$ Since b is arbitrary, UG: $\forall x(D(x) \rightarrow \neg S(x))$, the required claim.

5. Prove or Disprove.

$$\begin{split} & [\exists x P(x) \to \forall x Q(x)] \to \forall x [P(x) \to Q(x)] \\ & \text{. The Claim is true and we shall present below a proof.} \\ & \text{By definition; } \exists x P(x) \to \forall x Q(x) \equiv \neg \exists x P(x) \lor \forall x Q(x) \\ & \text{It follows that, } \forall x \neg P(x) \lor \forall x Q(x) \equiv \forall x (\neg P(x) \lor Q(x)) \\ & \text{Thus, } \forall x (P(x) \to Q(x)). \end{split}$$

6. $\forall x[P(x) \to Q(x)] \to [\exists x P(x) \to \forall x Q(x)]$

The above claim is false. We need a counter example which satisfies the premise and does not satisfy the conclusion. In particular, the example is such that $\exists x P(x)$ is true and $\forall x Q(x)$ is false. UOD: set of natural numbers, P(x) x is divisible by 6, Q(x) x is divisible by 2. It is clear that if a number is div by 6 then it is div by 2 and thus, $\forall x (P(x) \to Q(x))$ is true always. Also, $\exists x P(x)$ is true, for example x = 6. Further, $\forall x Q(x)$ is false, not every natural number is divisible by 2. Therefore, the conclusion is false.

7. Express the following using the first order logic by clearly mentioning UOD, predicates used: Everyone who gets admitted into an IIT gets a job. Therefore, if there are no jobs, then nobody gets admitted into any IIT.

Premise: $\forall x (person(x) \land (\exists y (IIT(y) \land admit(x, y))) \rightarrow (\exists z (job(z) \land getjob(x, z))))$ Conclusion: $\forall z (\neg (job(z)) \rightarrow \neg (\exists x (person(x) \land \exists y (IIT(y) \land admit(x, y)))).$ (or) $\neg \exists z ((job(z)) \rightarrow \forall x (person(x) \land \exists y (IIT(y)) \rightarrow \neg admit(x, y)).$

The claim is false. Consider a venn diagram with two sets A: set of jobs B: students admitted in IITs. $A \wedge B$ denotes, students of some IIT gets job. $B \setminus A$ denotes students in other IITs.